

CS 331, Fall 2025
Lecture 18 (10/29)

Today: - PSD matrices
- Convexity in $\mathbb{R}^d$
- Linear regression
- Algos for LP

PSD matrices (Part VI, Section 3.3)

SVD $A = U \Sigma V^T$, $\Sigma = \text{diag}(\sigma)$
nonnegative

$$A = \sum_{i \in [d]} \sigma_i u_i v_i^T$$

Column of U Column of V

Interpretation: each unit of $v_i \rightarrow \sigma_i$ units of u_i

$$\underbrace{Ax}_{= Vc = \sum_{i \in [d]} c_i v_i} = U \underbrace{\Sigma V^T V}_{\text{orthonormal}} c = U \Sigma c = \sum_{i \in [d]} \sigma_i c_i u_i$$

Spectral theorem: $\forall$ symmetric $M \in \mathbb{R}^{d \times d}$

$$M = U \Lambda U^T$$

$$U^T U = I$$

"orthonormal"

$$= \underbrace{\sum_{i \in [d]} \lambda_i u_i u_i^T}_{\text{eigendecomposition}}$$

$$U = \begin{pmatrix} u_1 & u_2 & \dots & u_d \end{pmatrix}$$

$$\Lambda = \begin{pmatrix} \lambda_1 & & 0 \\ & \lambda_2 & \\ 0 & & \ddots \\ & & & \lambda_d \end{pmatrix}$$

Compare to SVD:

$$M = U \Sigma V^T$$

$$= \sum_{i \in [d]} \sigma_i u_i v_i^T$$

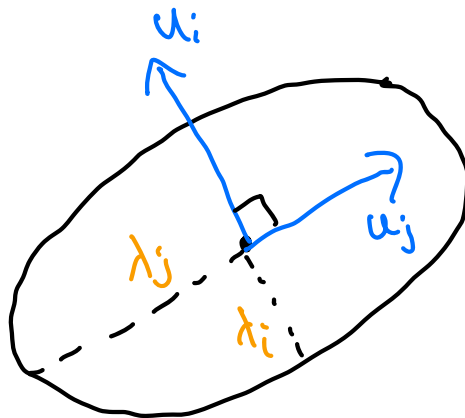
idea: if $\lambda_i < 0$

$$\Rightarrow \sigma_i = -\lambda_i$$
$$v_i = -u_i$$

If all $\lambda_i \geq 0$ also, SVD = eigendecomposition.

We call such M PSD

"Every PSD matrix is an ellipse"



We say (u, λ) is eigvec / eigenval of M
 if: $Mu = \lambda u$ "eigpair"

Claim: if $M = U \Lambda U^T = \sum_{i \in [d]} \lambda_i u_i u_i^T$

then (λ_j, u_j) is eigpair $\forall j \in [d]$

Proof: $\left(\sum_{i \in [d]} \lambda_i u_i u_i^T \right) u_j = \lambda_j u_j$

$$\lambda_1 \underbrace{u_1 u_1^T u_j}_{=0} + \dots + \lambda_j \underbrace{u_j u_j^T u_j}_{=1} + \dots$$

Main claim: let M be symmetric, s.d.

It's PSD iff $v^T M v \geq 0 \quad \forall v \in \mathbb{R}^d$

Proof: let $M = \sum_{i \in [d]} \lambda_i u_i u_i^T$

M not PSD $\Rightarrow \lambda_i < 0$ for some i

$$\Rightarrow u_i^T M u_i = u_i^T (\lambda_i u_i) < 0$$

$$M \text{ is PSD} \Rightarrow v^T \left(\sum_{i \in [d]} \lambda_i u_i u_i^T \right) v$$

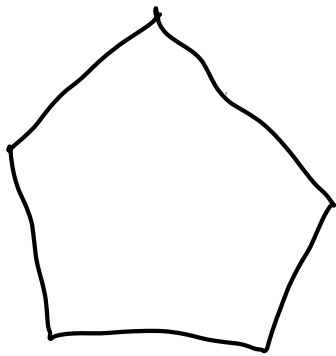
$$= \sum_{i \in [d]} \lambda_i \underbrace{(u_i^T v)}_{\geq 0}^2 \geq 0.$$

e.g. M PSD means $M_{ii} \geq 0$: let $v = e_i$

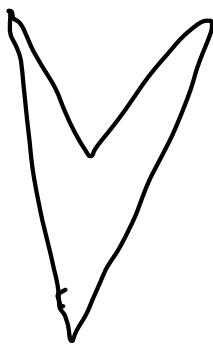
Convexity in $\mathbb{R}^d$ (Part VI, Section 5.1)

Let $X \subseteq \mathbb{R}^d$. We say X is convex

if it contains all lines:



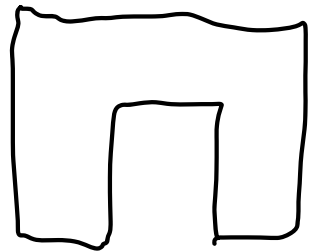
✓



✗



✓



✗

We say $f: X \rightarrow \mathbb{R}$ is convex if:
 $X \subseteq \mathbb{R}^d$

- X convex

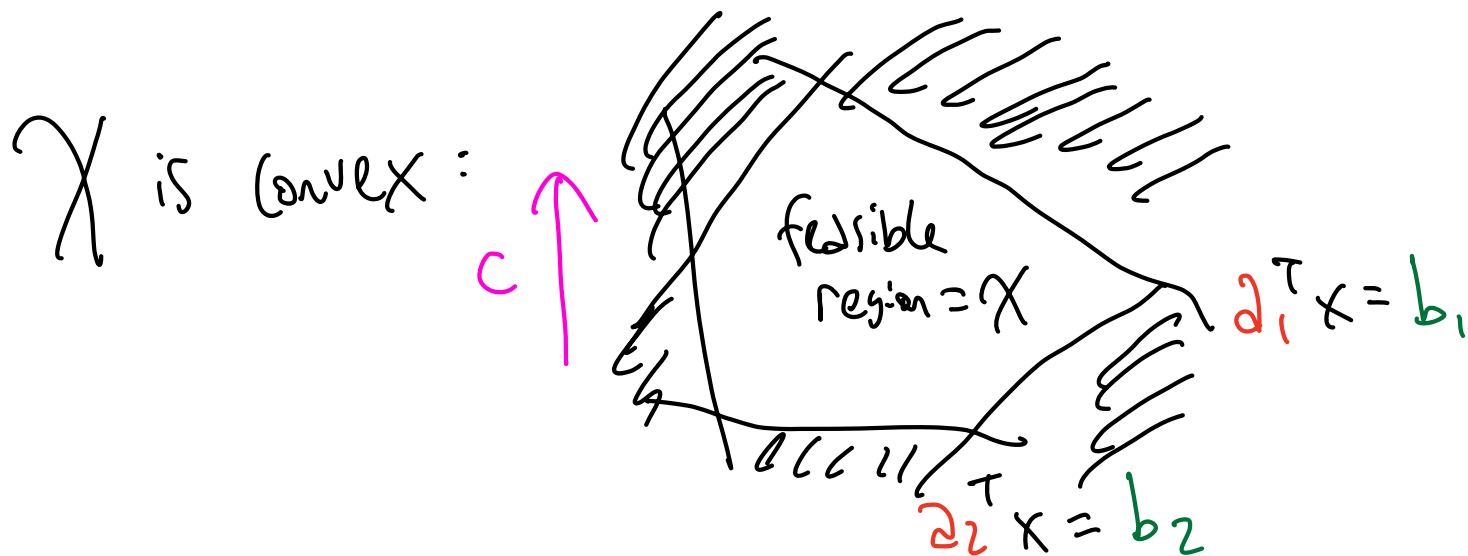
- All 1-D restrictions of f are convex
(stay at/below the line)

Two examples so far:

Section 3 • $X = \mathbb{R}$ • $d=1$

Section 2 • $X = \{x \in \mathbb{R}^d \mid Ax \leq b\}$

• $f(x) = c^T x$ (LP)



If $a_i^T x \leq b_i \quad \forall i \in [d],$

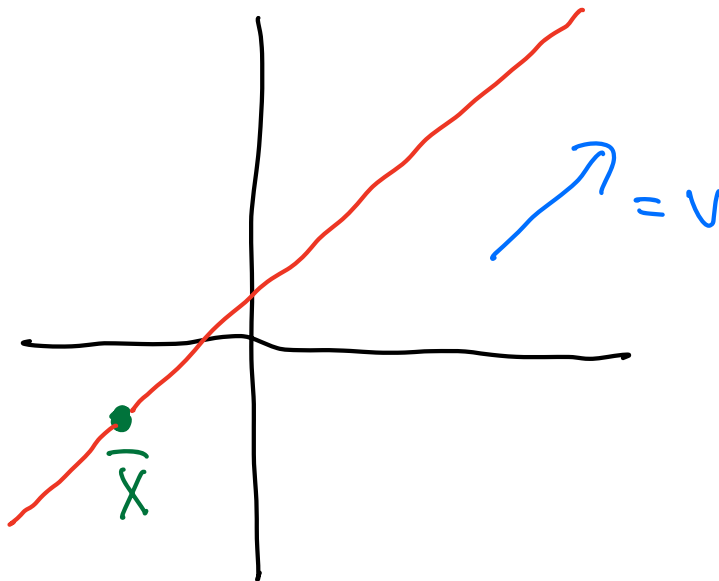
$a_i^T y \leq b_i$

$a_i^T \left(\underbrace{(1-\lambda)x + \lambda y}_{\text{line between } x \text{ \& } y} \right) \leq b_i \quad \forall \lambda \in [0,1]$

line between x & y

Let $f: X \rightarrow \mathbb{R}$ be convex, so:

$$f_0(\textcolor{red}{+}) \\ = f(\textcolor{green}{\bar{x}} + \textcolor{red}{+} \textcolor{blue}{v}) \\ \text{is convex.}$$



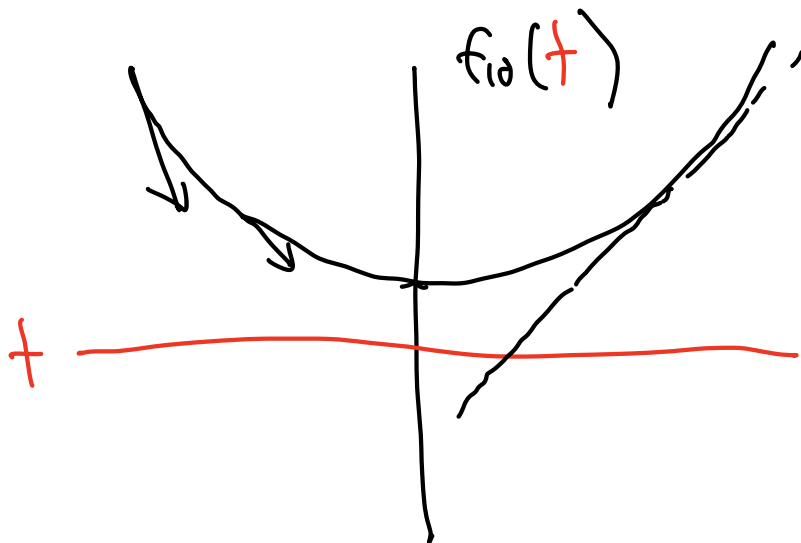
Fact 1:

$$f_0(\textcolor{red}{u}) \geq f_0(\textcolor{red}{+}) + f'_0(\textcolor{red}{+})(\textcolor{red}{u} - \textcolor{red}{+})$$

"above the tangent line"

Fact 2:

$$f''_0(\textcolor{red}{+}) \geq 0$$



So, what is $f'_{10}(+)$?

$$\frac{\partial}{\partial +} f(\underbrace{\bar{x} + + v}_x) = \frac{\partial f}{\partial x_1}(x) \frac{\partial x_1}{\partial +} + \dots$$

$$= \frac{\partial f}{\partial x_1}(x) v_1 + \dots + \frac{\partial f}{\partial x_d}(x) v_d = \nabla f(x)^T v$$

$$\text{Similarly, } f''_{10}(+) = \underbrace{v^T \nabla^2 f(x) v}_{d \times d \text{ matrix of } \frac{\partial^2 f}{\partial x_i \partial x_j}}$$

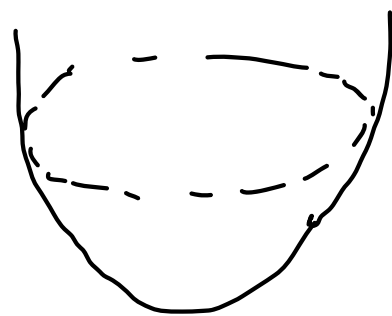
In $\mathbb{R}^d$: f is convex

$$\text{Fact 1: } f(y) \geq f(x) + \underbrace{\nabla f(x)^T (y-x)}_{\text{the "tangent line"}}$$

$$\text{Fact 2: } v^T \nabla^2 f(x) v \geq 0 \quad \forall v \in \mathbb{R}^d$$

Punchline: in $\mathbb{R}^D$, f is convex

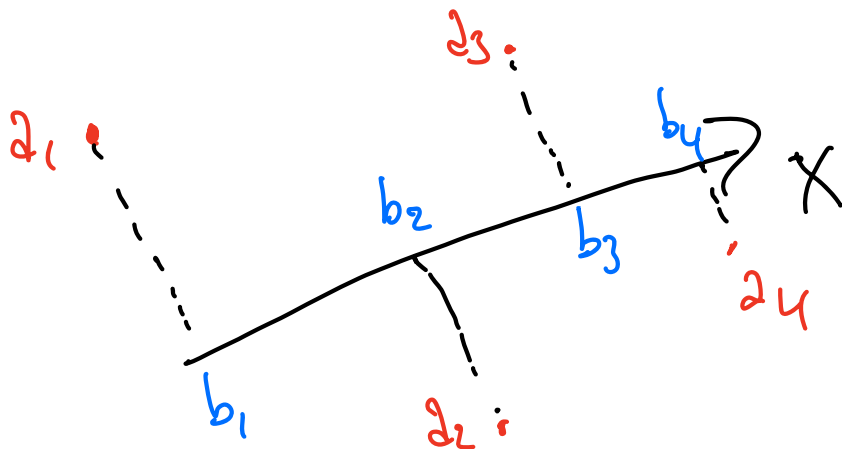
$\Leftrightarrow \nabla^2 f(x)$ is PSD



Linear regression (Part VI, Section 5.2)

Input: $\left\{ \underbrace{\left(\underset{\substack{\uparrow \\ \mathbb{R}^D}}{a_i}, \underset{\substack{\uparrow \\ \mathbb{R}}}{b_i}} \right)}_{\text{"feature vector"}} \dots \left(\underset{\substack{\uparrow \\ \mathbb{R}^D}}{a_n}, \underset{\substack{\uparrow \\ \mathbb{R}}}{b_n} \right) \right\}$
"response"

Output: X that "explains" relationship



$$\text{Goal: } \min_{x \in \mathbb{R}^d} \sum_{i \in [n]} \left(\underset{\nearrow}{a_i^T} x - b_i \right)^2$$

Model of relationship is linear
(HW: piecewise const.)

$$\text{Another way to write: } A = \begin{pmatrix} a_1^T \\ a_2^T \\ \vdots \\ a_n^T \end{pmatrix}$$

$$f(x) = \|Ax - b\|_2^2$$

$$= (Ax - b)^T (Ax - b)$$

$$= x^T \underbrace{A^T A}_M x - 2b^T A x + \text{const.}$$

Recall: M is PSD.

Fact: if $f(x) = x^T M x + v^T x$
then $\nabla f(x) = 2M$

Regression is convex ☺

For "nice" M we can solve fast using

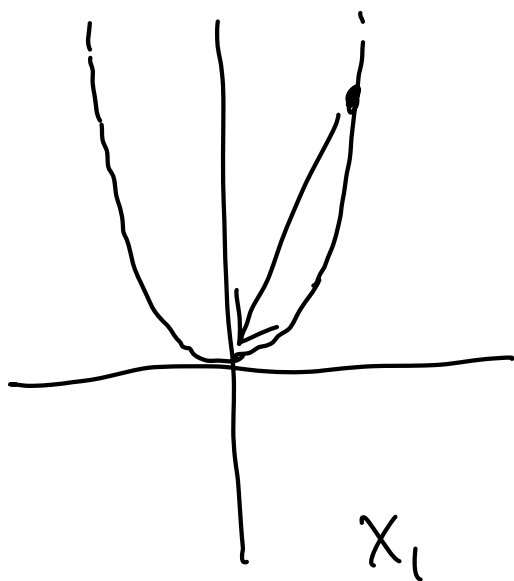
Gradient descent

Sketch: Suppose $M = \begin{pmatrix} \lambda_1 & 0 \\ 0 & \lambda_2 \end{pmatrix}$ (diagonal)

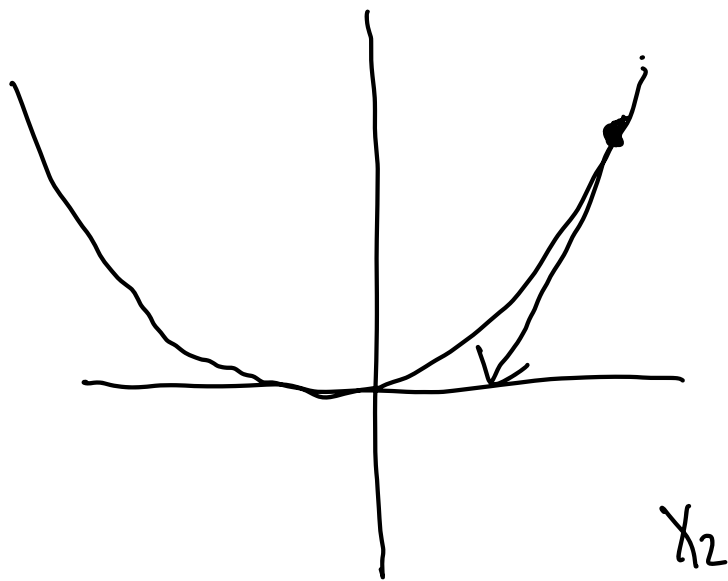
$$f(x) = \left(\lambda_1 x_1^2 + v_1 x_1 \right) + \left(\lambda_2 x_2^2 + v_2 x_2 \right)$$

only get to pick one step size...

$$\eta = \frac{1}{2\lambda_1} \text{ perfect for } x_1$$



perfect



undershot by factor $\frac{\lambda_2}{\lambda_1}$

Takes $\approx \frac{\lambda_1}{\lambda_2}$ iterations to converge.

General case: $M = U \begin{pmatrix} \lambda_1 & & \\ & \lambda_2 & \\ & & \ddots \\ & & & \lambda_d \end{pmatrix} U^T$

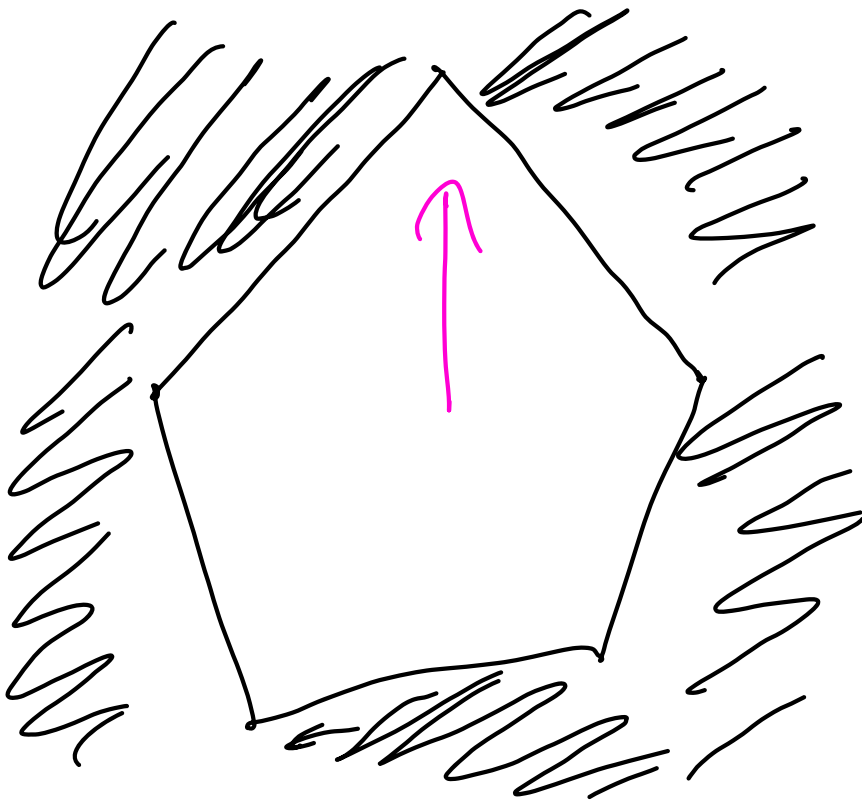
- GD converges in $\approx \frac{\lambda_1}{\lambda_d}$ iterations

"how spherical is M "?

- No problem if $U \neq I$

"input world = output world"

Algos for LP (Part VI, Section 5.3)



It's convex, but...

- Constrained
- Non-Smooth

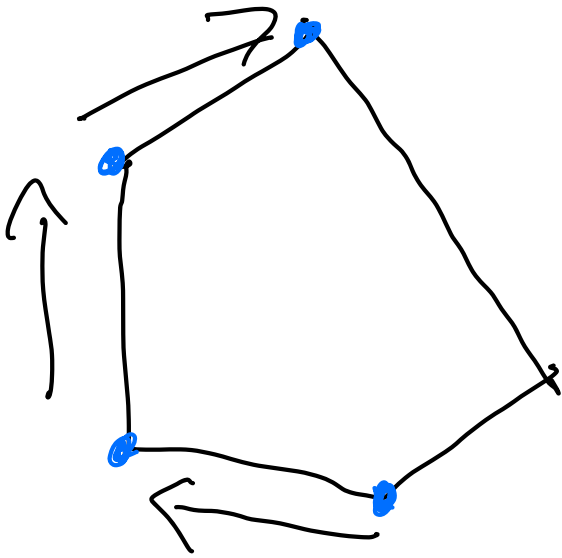
how to solve?

Strongly polynomial: Open "Smale's 9th problem"

Weakly polynomial: ✓ Great in practice!

- 10 most influential of 20th Century
- NYT headline: "Breakthrough in Problem Solving"

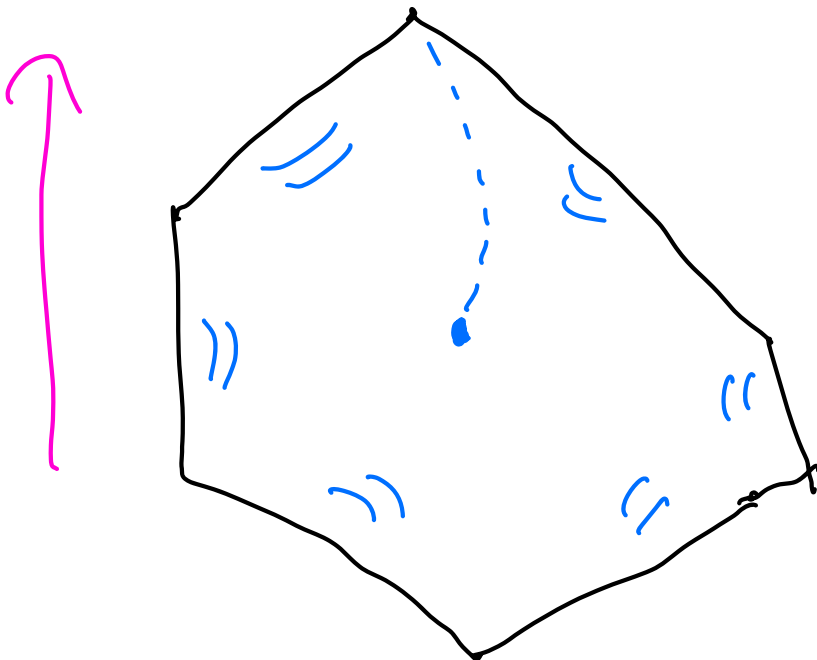
Idea 1: Simplex "walk along the boundary"



- decent in practice
- worst-case exponential
- "most instances" polynomial

Idea 2: IPM "decrease the force fields"

$$x_+ = \arg\max C^T x + \underbrace{t b(x)}_{\text{"barrier"}}$$



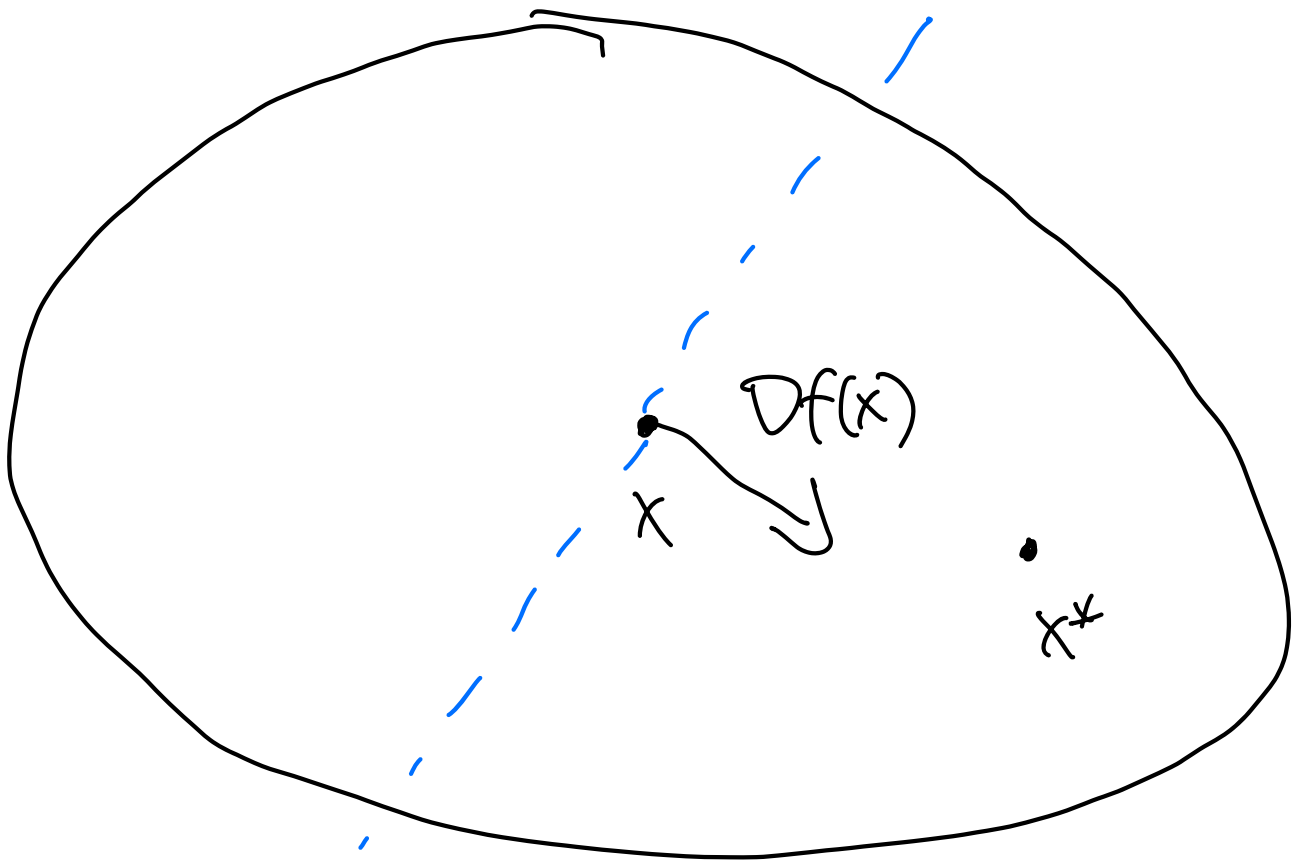
- polynomial
- SOTA in theory
- SOTA (mostly) in practice

Idea 3: CPM "binary search in $\mathbb{R}^d$ "

- Works for all convex functions
- My favorite algo 😊

In 1-D: $f'(x)$ lets you binary search

In $\mathbb{R}^d$: Use $\nabla f(x)$



If planes through center-of-gravity x ,
 $\geq 30\%$ on both sides. Decrease volume!